Order of convergence

Suppose $\{p_n\}_{n=0}^{\infty}$ is a sequence that converges to p, with $p_n \neq p$ for all n. If positive constants λ and α exist with

$$\lim_{n\to\infty}\frac{|p_{n+1}-p|}{|p_n-p|^{\alpha}}=\lambda,$$

then $\{p_n\}_{n=0}^{\infty}$ converges to p of order α , with asymptotic error constant λ .



Linear and Quadratic Order of convergence

- (i) If $\alpha = 1$ (and $\lambda < 1$), the sequence is **linearly convergent**.
- (ii) If $\alpha = 2$, the sequence is quadratically convergent.

Computing square root with Newton's Method

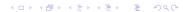
▶ Given a > 0, $p \stackrel{def}{=} \sqrt{a}$ is positive root of equation

$$f(x) \stackrel{\text{def}}{=} x^2 - a = 0.$$

Newton's Method

$$p_{k+1} = p_k - \frac{p_k^2 - a}{2p_k} = \frac{1}{2} \left(p_k + \frac{a}{p_k} \right), k = 0, 1, 2, \cdots,$$

▶ Newton's Method is well defined for any $p_0 > 0$.



Newton's Method for \sqrt{a} converges quadratically if $p_0 > 0$

• $\{p_k\}_{k=1}^{\infty}$ bounded below by \sqrt{a} :

$$p_k - \sqrt{a} = \frac{1}{2} \left(p_{k-1} + \frac{a}{p_{k-1}} \right) - \sqrt{a} = \frac{\left(p_{k-1} - \sqrt{a} \right)^2}{2p_{k-1}} \ge 0, \quad k = 1, 2, \cdots,$$

• $\{p_k\}_{k=1}^{\infty}$ monotonically decreasing:

$$p_{k+1} - p_k = \frac{1}{2} \left(p_k + \frac{a}{p_k} \right) - p_k = \frac{a - p_k^2}{2p_k} \le 0, \quad k = 1, 2, \dots,$$

▶ $\lim_{k\to\infty} p_k \stackrel{\text{def}}{=} p$ exists and satisfies Newton Iteration:

$$p=rac{1}{2}\left(p+rac{a}{p}
ight)\geq\sqrt{a}, \quad ext{therefore} \quad p=\sqrt{a}.$$

Newton's Method quadratically convergent

$$\lim_{k \to \infty} \frac{|p_k - \sqrt{a}|}{|p_{k-1} - \sqrt{a}|^2} = \lim_{k \to \infty} \frac{1}{2p_{k-1}} = \frac{1}{2\sqrt{a}}.$$

Linear Convergence Theorem of Fixed Point Iteration

- ▶ Let $g \in C[a, b]$ satisfy $g(x) \in [a, b]$ for all $x \in [a, b]$.
- ▶ **Assume** g'(x) continuous with $|g'(x)| \le \kappa$ for all $x \in (a, b)$ and constant $\kappa < 1$.
- ▶ **Then** Fixed Point Iteration

$$p_k=g(p_{k-1}), k=1,2,\cdots,$$

for any $p_0 \in [a,b]$ converges to fixed point $p \in [a,b]$ with

$$\lim_{k\to\infty}\frac{|p_{k+1}-p|}{|p_k-p|}=|g'(p)|\leq\kappa,\quad\text{if }p_k\neq p\text{ for all }k.$$

(not linear convergence if g'(p) = 0)

Proof of Linear Convergence

- ▶ By Fixed Point Theorem, Fixed Point $p \in [a, b]$ uniquely exists, and FPI converges to p.
- ▶ By Mean Value Theorem, there exists a ξ_k in between p_k and p such that

$$g(p_k)-g(p)=g'(\xi_k)(p_k-p),$$
 for all k .

Therefore

$$\lim_{k \to \infty} g'(\xi_k) = g'(p)$$

$$\lim_{k \to \infty} \frac{|p_{k+1} - p|}{|p_k - p|} = \lim_{k \to \infty} \frac{|g(p_k) - g(p)|}{|p_k - p|} = \lim_{k \to \infty} \frac{|g'(\xi_k)(p_k - p)|}{|p_k - p|}$$

$$= |g'(p)|$$

Quadratic Convergence of Fixed Point Iteration

▶ **Assume** FPI $\{p_k\}_{k=1}^{\infty}$ converges to p, with g'(p) = 0, then

$$\lim_{k\to\infty}\frac{|p_{k+1}-p|}{\left|p_k-p\right|^2}=\frac{1}{2}|g''(p)|,\quad \text{if }p_k\neq p\text{ for all }k.$$

(not quadratic convergence if g''(p) = 0)

Proof of Quadratic Convergence

Taylor expansion at p:

$$g(p_k) = \frac{\xi_k \text{ between } p, \ p_k}{g(p) + g'(p)(p_k - p) + \frac{1}{2}g''(\xi_k)(p_k - p)^2}$$

$$= g(p) + \frac{1}{2}g''(\xi_k)(p_k - p)^2$$

Therefore

$$\lim_{k \to \infty} \frac{|p_{k+1} - p|}{|p_k - p|^2} = \lim_{k \to \infty} \frac{|g(p_k) - p|}{|p_k - p|^2} = \frac{1}{2} \lim_{k \to \infty} \frac{|g''(\xi_k)(p_k - p)^2|}{|p_k - p|^2}$$
$$= \frac{1}{2} |g''(p)|$$

Quadratic Convergence of Newton Iteration

Newton Iteration is FPI with

$$g(x) = x - \frac{f(x)}{f'(x)},$$

and with f(p) = 0.

derivatives

$$g'(x) = \frac{f(x)f''(x)}{(f'(x))^2}$$

$$g''(x) = \frac{f''(x)}{f'(x)} + \frac{f'(x)f'''(x) - 2(f''(x))^2}{(f'(x))^3} f(x)$$

therefore

$$g'(p) = 0$$
, and $g''(p) = \frac{f''(p)}{f'(p)}$,
 $\lim_{k \to \infty} \frac{|p_{k+1} - p|}{|p_k - p|^2} = \frac{1}{2} \left| \frac{f''(p)}{f'(p)} \right|$, if $f'(p) \neq 0$.



Multiple roots:
$$f(p) = 0, f'(p) = 0$$

▶ **Definition**: A solution p of f(x) = 0 is a root of multiplicity m of f if for $x \neq p$, we can write

$$f(x) = (x - p)^m q(x)$$
, where $\lim_{x \to p} q(x) \neq 0$.

▶ **Theorem**: The function $f \in C^m[a, b]$ has a root of multiplicity m at p in (a, b) if and only if

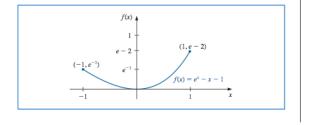
$$0 = f(p) = f'(p) = f''(p) = \cdots = f^{(m-1)}(p)$$
, but $f^{(m)}(p) \neq 0$.

simple root: f satisfies f(p) = 0, but $f'(p) \neq 0$.



Example:
$$f(x) = e^x - x - 1$$
, $f(0) = f'(0) = 0$, $f''(0) = 1$

Function

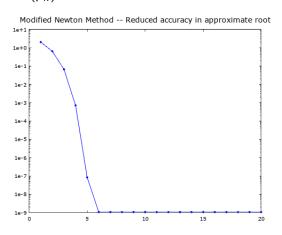


Newton's Method

n	p_n		
0	1.0		
1	0.58198		
2	0.31906		
3	0.16800		
4	0.08635		
5	0.04380		
6	0.02206		
7	0.01107		
8	0.005545		
9	2.7750×10^{-3}		
10	1.3881×10^{-3}		
11	6.9411×10^{-4}		
12	3.4703×10^{-4}		
13	1.7416×10^{-4}		
14	8.8041×10^{-5}		
15	4.2610×10^{-5}		
16	1.9142×10^{-6}		

Modified Newton Method

$$p_{k+1} = p_k - \frac{m f(p_k)}{f'(p_k)}$$
, for $k = 1, 2, \dots$, $(m = \text{multiplicity of root } p.)$



•
$$m=2$$
 for function $f(x)=e^x-x-1$,



Modified Newton Method is Quadratically Convergent

$$p_{k+1} = g(p_k), \quad g(x) = x - \frac{m f(x)}{f'(x)}, \text{ for } k = 1, 2, \dots,$$

▶ Let $f(x) = (x - p)^m q(x)$, $q(p) \neq 0$.

•

$$g(x) = x - \frac{m f(x)}{f'(x)} = x - \frac{m (x - p)q(x)}{mq(x) + (x - p)q'(x)}.$$

• g'(p) = 0, hence quadratic convergence.

Accelerating Convergence: Aitken's Δ^2 Method

▶ **Suppose** $\{p_k\}_{k=1}^{\infty}$ linearly converges to limit p,

$$\lim_{k\to\infty}\frac{p_{k+1}-p}{p_k-p}=\lambda,\quad |\lambda|<1.$$

Define

$$\frac{p_{k+1}-p}{p_k-p}\stackrel{def}{=} \lambda_k,$$

so that $\{\lambda_k\}_{k=1}^{\infty}$ converges to λ .

▶ It follows that

$$0 \approx \lambda_{k+1} - \lambda_k = \frac{p_{k+2} - p}{p_{k+1} - p} - \frac{p_{k+1} - p}{p_k - p}.$$

► Solve for *p*:

$$p = \frac{p_{k+1}^2 - p_k p_{k+2}}{2p_{k+1} - p_k - p_{k+2}} + \frac{(\lambda_{k+1} - \lambda_k)(p_{k+1} - p)(p_k - p)}{2(p_{k+1} - p) - (p_k - p) - (p_{k+2} - p)}.$$

Accelerating Convergence: Aitken's Δ^2 Method

$$\widehat{p}_{k} \stackrel{\text{def}}{=} \frac{p_{k+1}^{2} - p_{k}p_{k+2}}{2p_{k+1} - p_{k} - p_{k+2}} \\
= p_{k} - \frac{(p_{k+1} - p_{k})^{2}}{p_{k+2} - 2p_{k+1} + p_{k}} \stackrel{\text{def}}{=} \{\Delta^{2}\}(p_{k}).$$

Approximation Error

$$|\widehat{p}_{k} - p| = \left| \frac{(\lambda_{k+1} - \lambda_{k})(p_{k+1} - p)(p_{k} - p)}{2(p_{k+1} - p) - (p_{k} - p) - (p_{k+2} - p)} \right|$$

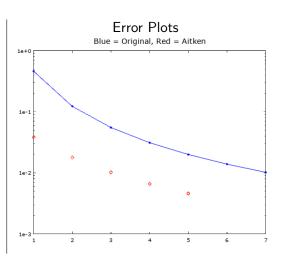
$$= \left| \frac{(\lambda_{k+1} - \lambda_{k})(p_{k} - p)}{2 - \left(\frac{p_{k+1} - p}{p_{k} - p}\right)^{-1} - \frac{p_{k+2} - p}{p_{k+1} - p}} \right|$$

$$\approx \left| \frac{(\lambda_{k+1} - \lambda_{k})(p_{k} - p)}{2 - \lambda^{-1} - \lambda} \right| \ll O(|p_{k} - p|).$$

Accelerating Convergence: Aitken's Δ^2 Method

p_n	=	cos((1/	n)
-------	---	------	-----	----

n	p_n	\hat{p}_n
1	0.54030	0.96178
2	0.87758	0.98213
3	0.94496	0.98979
4	0.96891	0.99342
5	0.98007	0.99541
6	0.98614	
7	0.98981	



Accelerating Convergence: Steffensen's Method

▶ Aitken's Acceleration for a given $\{p_k\}_{k=1}^{\infty}$:

$$\{\Delta^2\}(p_k) = p_k - \frac{(p_{k+1} - p_k)^2}{p_{k+2} - 2p_{k+1} + p_k}.$$

- ► Steffensen's Method: use Aitken's Acceleration as soon as iterations are generated:
 - Given $p_0^{(0)}$
 - for $k = 0, 1, 2, \cdots$,

$$p_1^{(k)} = g(p_0^{(k)}), \ p_2^{(k)} = g(p_1^{(k)}), \ p_0^{(k+1)} = \{\Delta^2\}(p_0^{(k)}).$$

Steffensen's

To find a solution to p = g(p) given an initial approximation p_0 :

INPUT initial approximation p_0 ; tolerance TOL; maximum number of iterations N_0 .

OUTPUT approximate solution p or message of failure.

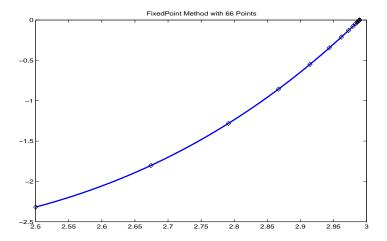
- Step 1 Set i = 1.
- Step 2 While $i \le N_0$ do Steps 3-6.

Step 3 Set
$$p_1 = g(p_0)$$
; (Compute $p_1^{(i-1)}$.)
$$p_2 = g(p_1)$$
; (Compute $p_2^{(i-1)}$.)
$$p = p_0 - (p_1 - p_0)^2 / (p_2 - 2p_1 + p_0)$$
. (Compute $p_0^{(i)}$.)

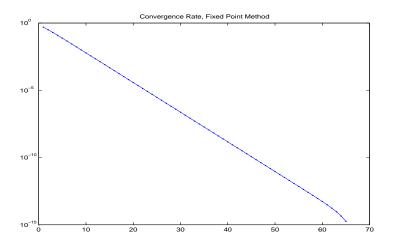
Step 4 If
$$|p-p_0| < TOL$$
 then OUTPUT (p) ; (Procedure completed successfully.) STOP.

- **Step 5** Set i = i + 1.
- **Step 6** Set $p_0 = p$. (Update p_0 .)
- Step 7 OUTPUT ('Method failed after N_0 iterations, $N_0 =$ ', N_0); (Procedure completed unsuccessfully.) STOP.

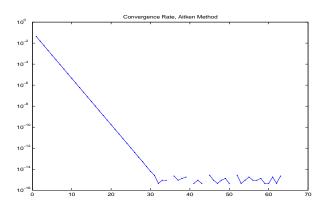
Fixed point for $g(x) = log(2 + 2x^2)$



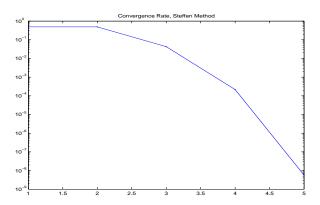
Fixed Point Iteration: Linear Convergence



Aitken's Method: Faster, but Linear Convergence



Steffenson's Method: Quadratic Convergence



Newton's Method on Polynomial roots = Horner's Method

Let

$$P(x) = a_n x_n + a_{n-1} x_{n-1} + \cdots + a_1 x + a_0.$$

Horner's Method computes $P(x_0)$ and $P'(x_0)$

- ▶ **define** $b_n = a_n$
- for $k = n 1, n 2, \dots, 1, 0$

$$b_k = a_k + b_{k+1} x_0,$$

▶ then

$$b_0 = P(x_0)$$

$$P(x) = (x - x_0) Q(x) + b_0$$

$$Q(x) = b_n x_{n-1} + b_{n-1} x_{n-2} + \dots + b_2 x + b_1.$$

▶ it follows

$$P'(x) = Q(x) + (x - x_0) Q'(x)$$

 $P'(x_0) = Q(x_0)$



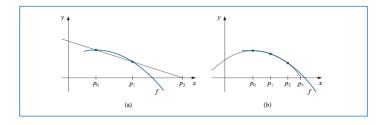
Horner's Method

```
INPUT degree n; coefficients a_0, a_1, \ldots, a_n; x_0.
OUTPUT y = P(x_0); z = P'(x_0).
Step 1 Set y = a_n; (Compute b_n for P.)
             z = a_n. (Compute b_{n-1} for O.)
Step 2 For i = n - 1, n - 2, \dots, 1
              set y = x_0 y + a_i; (Compute b_i for P.)
                 z = x_0 z + y. (Compute b_{i-1} for Q.)
Step 3 Set y = x_0y + a_0. (Compute b_0 for P.)
Step 4 OUTPUT (y,z);
          STOP
```

Muller's Method: finding complex roots

- ▶ **Given** three points $(p_0, f(p_0)), (p_1, f(p_1)),$ and $(p_2, f(p_2)).$
- Construct a parabola through them,
- \triangleright p_3 is the intersection of x-axis with parabola closer to p_2 .

Secant Method vs. Muller's Method



Muller's Method: derivation

Choose parabola

$$P(x) = a(x - p_2)^2 + b(x - p_2) + c.$$

a, b, c satisfy

$$f(p_0) = a(p_0 - p_2)^2 + b(p_0 - p_2) + c,$$

$$f(p_1) = a(p_1 - p_2)^2 + b(p_1 - p_2) + c,$$

$$f(p_2) = a \cdot 0^2 + b \cdot 0 + c = c$$

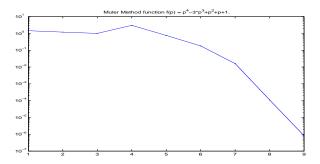


Muller's Method: finding complex roots

 \triangleright p_3 is the intersection of x-axis with parabola closer to p_2 .

$$p_3 = p_2 - \frac{2c}{b + \operatorname{sgn}(b)\sqrt{b^2 - 4ac}},$$

Muller's Method



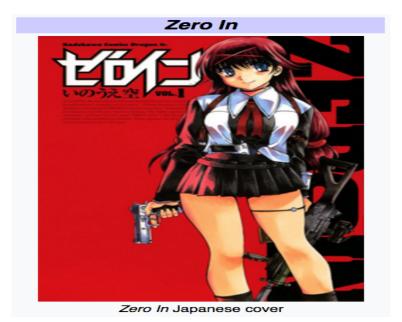
Choices, and more Choices

- ▶ Bisection Method: slow (linearly convergent); reliable (always finds a root given interval [a, b] with $f(a) \cdot f(b) < 0$.)
- ▶ Fixed Point Iteration: slow (linearly convergent); no need for interval [a, b] with $f(a) \cdot f(b) < 0$.
- Newton's Method: fast (quadratically convergent); could get burnt (need not converge.)
- Secant Method: between Bisection and Newton in speed; need not converge; no derivative needed.
- Steffensen's Method: fast (quadratically convergent); no derivative needed. not a method of choice in higher dimensions.
- Muller's Method: can handle complex roots; need not converge.

Brent's Method: Motivation

- ► (Mostly) speed of Steffensen's Method
- Reliability of Bisection Method

Brent's Method = Zeroin



Brent's Method: Zeroin (Wikipedia)

Brent's method

From Wikipedia, the free encyclopedia

In numerical analysis, **Brent's method** is a root-finding algorithm combining the bisection method, the secant method and inverse quadratic interpolation. It has the reliability of bisection but it can be as quick as some of the less-reliable methods. The algorithm

Brent's Method: Zeroin



Cleve's Corner: Cleve Moler on Mathematics and Computing

Scientific computing, math & more

< Zeroin, Part 1: Dekker's Algorithm | Zeroin, Part 3: MATLAB Zero.

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Zeroin, Part 2: Brent's Version

Posted by Cleve Moler, October 26, 2015

Richard Brent's improvements to Dekker's zeroin algorithm, published in 1971, made it faster, safer in floating point arithmetic, and quaranteed not to fail.

Contents

- · Richard Brent
- · Weakness of Zeroin
- · Two improvements
- · Muller's method
- · Inverse quadratic interpolation
- · Brent's algorithm
- Fzero



Brent's Method

▶ Worst case number of iterations = (n^2) , where n is number of iterations of Bisection Method for a given tolerance.

Finding Zeros of Single-Variable, Real Functions

Gautam Wilkins
University of California, San Diego

Modified Brent's Method

Modified Zero-In

- Brent's Method may be modified to ensure O(n) time instead of $O(n^2)$.
- Force a bisection if:
 - 1) If the size of the original interval is not reduced by a factor of 1/2 after five interpolation steps.
 - 2) If an interpolation step produces a point, x, such that |f(x)| is not a factor of 1/2 smaller than the previous best point.

Modified Brent's Method

Modified Zero-In

- The first condition ensures that the complexity of the search is O(n).

 The second condition addresses the issue of very flat functions, for which Brent's Method converges rather slowly.

Modified Brent's Method

Comparison

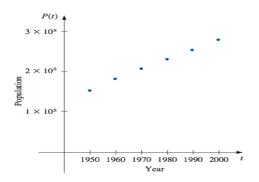
 For the worst-case function shown earlier, when Brent's Method took 2914 iterations, Modified Zero-In took 85 iterations.

 This reduction in complexity, as far as we can tell, comes at virtually no cost to performance in general.

Interpolation and Approximation (connecting the dots)

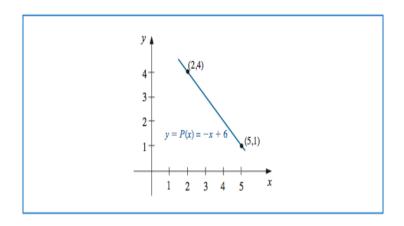
US population census data available every ten years.

Year	1950	1960	1970	1980	1990	2000
Population (in thousands)	151,326	179,323	203,302	226,542	249,633	281,422



- ▶ What was US population in year 1996?
- What will US population be in year 2017?

Connecting two dots



► Given two distinct dots, there is a unique line connecting them.

Connecting two dots

▶ Given two distinct points $(p_0, f(p_0))$ and $(p_1, f(p_1))$, there is a unique line connecting them

$$P(x) = f(p_0) \cdot \frac{x - p_1}{p_0 - p_1} + f(p_1) \cdot \frac{x - p_0}{p_1 - p_0}$$

$$\stackrel{def}{=} f(p_0) \cdot L_0(x) + f(p_1) \cdot L_1(x),$$

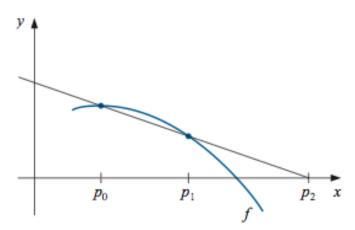
with $L_0(x), L_1(x)$ unrelated to $f(p_0), f(p_1)$.

$$L_0(p_0) = 1,$$
 $L_0(p_1) = 0,$
 $L_1(p_0) = 0,$ $L_1(p_1) = 1.$

Connecting two dots

▶ Given two distinct points $(p_0, f(p_0))$ and $(p_1, f(p_1))$, there is a unique line connecting them

$$P(x) = f(p_0) \cdot \frac{x - p_1}{p_0 - p_1} + f(p_1) \cdot \frac{x - p_0}{p_1 - p_0}$$



Connecting n + 1 dots

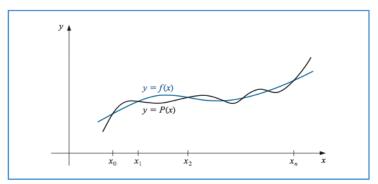
• Given n+1 distinct points

$$(x_0, f(x_0)), (x_1, f(x_1)), \cdots, (x_n, f(x_n)),$$

▶ Find a degree $\leq n$ polynomial

$$P(x) = a_n x^n + a_{n-1} x^{n-1} + \dots + a_1 x + a_0,$$

so that $P(x_0) = f(x_0), P(x_1) = f(x_1), \dots, P(x_n) = f(x_n).$



Connecting n + 1 dots

 \blacktriangleright Would like to write P(x) as

$$P(x) = f(x_0)L_0(x) + f(x_1)L_1(x) + \cdots + f(x_n)L_n(x),$$

where $L_0(x), L_1(x), \dots, L_n(x)$ are unrelated to $f(x_0), f(x_1), \dots, f(x_n)$.

▶ at each node x_j , $j = 0, 1, \dots, n$,

$$f(x_j) = P(x_j) = f(x_0)L_0(x_j) + f(x_1)L_1(x_j) + \cdots + f(x_n)L_n(x_j)$$

▶ This suggests for all i, j,

$$L_i(x_j) = \begin{cases} 1, & \text{if } i = j, \\ 0, & \text{if } i \neq j, \end{cases}$$



Largrangian Polynomials

$$P(x) = f(x_0)L_0(x) + f(x_1)L_1(x) + \cdots + f(x_n)L_n(x),$$

$$L_i(x_j) = \begin{cases} 1, & \text{if } i = j, \\ 0, & \text{if } i \neq j, \end{cases}$$

 x_j is a root of $L_i(x)$ for every $j \neq i$, so

$$L_i(x) = \prod_{i \neq i} \frac{x - x_j}{x_i - x_j}.$$

Quadratic Interpolation

- (a) Use the numbers (called *nodes*) $x_0 = 2$, $x_1 = 2.75$, and $x_2 = 4$ to find the second Lagrange interpolating polynomial for f(x) = 1/x.
- (b) Use this polynomial to approximate f(3) = 1/3.

$$L_0(x) = \frac{(x - 2.75)(x - 4)}{(2 - 2.5)(2 - 4)} = \frac{2}{3}(x - 2.75)(x - 4),$$

$$L_1(x) = \frac{(x - 2)(x - 4)}{(2.75 - 2)(2.75 - 4)} = -\frac{16}{15}(x - 2)(x - 4),$$

$$L_2(x) = \frac{(x-2)(x-2.75)}{(4-2)(4-2.5)} = \frac{2}{5}(x-2)(x-2.75).$$

$$P(x) = \sum_{k=0}^{2} f(x_k) L_k(x)$$

$$= \frac{1}{3} (x - 2.75)(x - 4) - \frac{64}{165} (x - 2)(x - 4) + \frac{1}{10} (x - 2)(x - 2.75)$$

$$= \frac{1}{22} x^2 - \frac{35}{88} x + \frac{49}{44}.$$

$$f(3) \approx P(3) = \frac{9}{22} - \frac{105}{88} + \frac{49}{44} = \frac{29}{88} \approx 0.32955.$$

